# Tutorial: <br> Determinantal Point Processes and their Application to Signal Processing and Machine Learning 

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Introduction
DPPs to produce diverse samples
DPPs as a tool in SP/ML
DPPs to characterize

Definition, basic properties
Repulsive point processes are hard DPPs, the nitty-gritty

Computation
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Conclusion

In a nutshell, determinantal point processes (or DPP) :

- are random processes that induce diversity.
- are tractable.
- are used for three main purposes:
i/ produce diverse samples of a large database
ii/ use as a tool in a variety of SP/ML contexts
iii/ characterize various observed phenomena.


## DPPs induce diversity



Figure: Example of iid uniform sampling

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Figure: Example of DPP sampling

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Figure: Example of DPP sampling
i/ This sample diversity can be directly useful ${ }^{12}$ :
summary generation:

${ }^{1}$ left figure: from Kulesza and Taskar, DPPs for machine learning, Found. and Trends in ML, 2013
${ }^{2}$ right figure: from G. Gautier's slides guilgautier.github.io/pdfs/GaBaVa17_slides.pdf
i/ This sample diversity can be directly useful ${ }^{12}$ :
summary generation:

search engines / recommendation:

'bole' query

[^0]i/ This sample diversity can be directly useful ${ }^{12}$ :
summary generation:

search engines / recommendation:

'bole' query
ii/ DPP samples can also be used as a tool in several SP/ML contexts:

- Monte Carlo integration
- Feature selection problems
- Coresets
- etc.

[^1]DPPs as a tool: an example


DPPs as a tool: an example


CPs as a tool: an example


Figure: Example of id uniform sampling

CPs as a tool: an example


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CPs as a tool: an example


Figure: Example of id uniform sampling

DPPs as a tool: an example


Figure: Example of iid uniform sampling

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CPs as a tool: an example


Figure: Example of id uniform sampling

CPs as a tool: an example


Figure: id estimations of the mean

DPPs as a tool: an example


Figure: Example of DPP sampling

DPPs as a tool: an example


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CPs as a tool: an example


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CPs as a tool: an example


Figure: Example of DPP sampling

CPs as a tool: an example


Figure: DPP estimations of the mean

DPPs as a tool: an example


Figure: Comparison of both estimators: variance reduction (here by a factor 3)
iii/ Finally, DPPs are used to characterize various phenomena.
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# Where do DPPs arise? 

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## Where do DPPs arise?



iii/ Finally, DPPs are used to characterize various phenomena.


## Where do DPPs arise?


-$-$

and more...

## Eigenvalues of the Gaussian Unitary Ensemble ${ }^{1}$

- Consider a Hermitian matrix $\mathrm{H} \in \mathbb{C}^{n \times n}$ with
- diagonal elements of the form $\mathrm{H}_{j j}=X$ with $X$ drawn iid from $\mathcal{N}(0,1)$
- off-diagonal elements of the form $\mathrm{H}_{j k}=X+i Y$ with $X$ and $Y$ drawn iid from $\mathcal{N}(0,1 / 2)$.

[^2]
## Eigenvalues of the Gaussian Unitary Ensemble ${ }^{1}$

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- off-diagonal elements of the form $\mathrm{H}_{j k}=X+i Y$ with $X$ and $Y$ drawn iid from $\mathcal{N}(0,1 / 2)$.
- It has $n$ real eigenvalues. They are distributed s.t.:

$$
\mathbb{P}\left(\lambda_{1}, \ldots, \lambda_{n}\right) \propto \exp ^{-\sum_{j} \lambda_{j}^{2}} \prod_{j<k}\left(\lambda_{j}-\lambda_{k}\right)^{2}
$$

[^3]
## Eigenvalues of the Gaussian Unitary Ensemble ${ }^{1}$

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- It has $n$ real eigenvalues. They are distributed s.t.:

$$
\begin{aligned}
\mathbb{P}\left(\lambda_{1}, \ldots, \lambda_{n}\right) & \propto \exp ^{-\sum_{j} \lambda_{j}^{2}} \prod_{j<k}\left(\lambda_{j}-\lambda_{k}\right)^{2} \\
& \propto \operatorname{det} \mathrm{M}^{2}
\end{aligned}
$$

where $M_{j k}=\lambda_{k}^{j-1} \exp ^{-\frac{1}{2} \lambda_{k}^{2}}$.

[^4]
## Eigenvalues of the GUE: illustration ${ }^{1}$

Examples of 6 point processes in 1D (3 GUE and 3 uniform):

${ }^{1}$ see, e.g., Johansson, Random matrices and DPPs, Arxiv (lecture notes), 2005

## Eigenvalues of the GUE: illustration ${ }^{1}$

Examples of 6 point processes in 1D (3 GUE and 3 uniform):


[^5]
## Eigenvalues of the GUE: illustration ${ }^{1}$

Examples of 6 point processes in 1D (3 GUE and 3 uniform):


[^6]
## A spinless fermion in a harmonic potential ${ }^{1}$



$$
V(x)=\frac{1}{2} x^{2}
$$

At temperature $T=0$, the probability distribution of the particle is a simple Gaussian:


[^7]
## Two non-interacting fermions ${ }^{1}$



Pauli's exclusion principle implies, after a few calculations, that, at $T=0$ :

$$
\begin{aligned}
\mathbb{P}\left(x_{1}, x_{2}\right) & \propto\left(x_{2}-x_{1}\right)^{2} e^{-\left(x_{1}^{2}+x_{2}^{2}\right)} \\
& \propto\left(\operatorname{det}\left[\begin{array}{cc}
e^{-\frac{1}{2} x_{1}^{2}} & e^{-\frac{1}{2} x_{2}^{2}} \\
x_{1} e^{-\frac{1}{2} x_{1}^{2}} & x_{2} e^{-\frac{1}{2} x_{2}^{2}}
\end{array}\right]\right)^{2}
\end{aligned}
$$

[^8]
## Two non-interacting fermions ${ }^{1}$

$$
\mathbb{P}\left(x_{2} \mid x_{1}=\bullet\right)
$$

Pauli's exclusion principle implies, after a few calculations, that, at $T=0$ :

$$
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\end{array}\right]\right)^{2}
\end{aligned}
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Interim: repulsive point processes are hard

- There are many ways of defining point processes that feature repulsion; some may look much more natural than DPPs
- An unfortunate fact of point process theory is that repulsive point processes are hard, theoretically and empirically
- Desirable features:

1. Probability density of p.p. is tractable (including normalisation constant)
2. Inclusion probabilities (intensity functions) are tractable
3. Sampling is tractable
4. Model is easy to understand

- DPPs have properties (1-3) and arguably (4) once you get used to them
- Most other repulsive processes have one or two (at best)


## Gibbs point processes

- Many repulsive point processes can be described using the general framework of Gibbs point processes
- A Gibbs point process takes the following form:

$$
p(\mathcal{X})=\frac{\exp \left(-\beta \sum_{i<j} v\left(x_{i}, x_{j}\right)\right)}{Z_{\beta}}
$$

- $v\left(x_{i}, x_{j}\right)$ is called a pairwise potential
- the sum runs over all pairs of points
- example : $v\left(x_{i}, x_{j}\right)=d\left(x_{i}, x_{j}\right)$ where $d$ is a distance, encourages points to be far apart.

The hard sphere model


The hard sphere model (AKA hard-core model) is used in physics to describe a set of particles that cannot overlap. See Löwen (2000) ${ }^{1}$.

[^10]The hard sphere model

- We assume that $\mathcal{X}=x_{1}, \ldots, x_{m}$, with $m$ fixed and $x_{i} \in[0,1]^{d}$
- The pairwise potential is simply:

$$
v\left(\boldsymbol{x}_{i}, \boldsymbol{x}_{j}\right)= \begin{cases}\infty & \text { if }\left\|\boldsymbol{x}_{i}-\boldsymbol{x}_{j}\right\|^{2}<r \\ 0 & \text { otherwise }\end{cases}
$$

## Things to think about

- What's the normalisation constant for the hard-sphere model? Hint: can you relate it to the probability that $m$ points sampled independently have a minimum pairwise distance $>r$ ?
- What are the valid configurations like when $m$ is large?
- How would you sample from the hard-sphere model?


## Normalisation constant

- Normalisation constant:

$$
\int_{\Omega^{m}} \prod_{i<j} \mathbb{I}\left(\left\|x_{i}-x_{j}\right\|^{2}>r\right) d x_{1} \ldots d x_{m}
$$

- Intractable (except in dimension one)!


## Packing limit

As $m$ becomes large, we reach the packing limit, and most configurations are impossible


In the general case packing is a very hard problem (image from Wikipedia)

## Sampling

- Possible sampling algorithm: "dart throwing".
- Pick a random initial location uniformly
- Pick a second location uniformly among remaining possible locations
- Pick a third location uniformly among remaining possible locations
- etc. until you have $m$ spheres or further sampling is impossible (start again)
- Very good for small $m$, very bard for large $m$


## Summary: the hard sphere model

- Simplest, most natural model you can imagine (property 4)
- But:

1. Probability density is intractable (because normalisation constant is intractable for $d>1$ )
2. Inclusion probabilities (intensity functions) are intractable for general domains, at least as far as we know
3. Sampling is easy for small $m$ (not very repulsive), then in large $m$ becomes equivalent to the notoriously hard sphere packing problem

## DPPs, the nitty-gritty

- We'll see that DPPs tick all boxes, contrary to most Gibbs processes
- The set-up cost is a bit higher; it's important to understand how these processes are defined, and to be careful about the notation
- We will now go through a few definitions in detail


## Some notation for discrete point processes

- $\Omega$ is a base set of size $n$ representing the items to sample from. w.l.o.g we may take $\Omega=\{1, \ldots, n\}$
- $\mathcal{X}$ is a random subset of $\Omega$
- We note $m=|\mathcal{X}|$, which may be a random variable


## L-ensembles

- The repulsion in DPPs is based on a notion of similarity between items in $\Omega$.
- The similarity between all pairs of items in $\Omega$ is stored in a $n \times n$ matrix called (for historical reasons) the "L-ensemble".
- We note this matrix $\mathbf{L}$, with $L_{i j}$ the similarity between items $i$ and $j$
- $\mathbf{L}$ is assumed to be positive definite.


## L-ensembles

- We'll come across several ways of constructing the $\mathbf{L}$ matrix.
- For now, assume that the items are vectors in $\mathbb{R}^{d}$. We can use a kernel function to describe similarity.
- Example: Gaussian kernel

$$
L_{i j}=\exp \left(-\frac{1}{2 \sigma^{2}}\left\|\mathbf{x}_{i}-\mathbf{x}_{j}\right\|^{2}\right)
$$

Similarity via the Gaussian kernel


Similarity via the Gaussian kernel


## DPP: formal definition

- We say that $\mathcal{X}$ (random set) is distributed according to a DPP if:

$$
p(\mathcal{X}=X) \propto \operatorname{det} \mathbf{L}_{X}
$$

- $\mathbf{L}_{\mathcal{X}}$ is the restriction of $\mathbf{L}$ to the items in $\mathcal{X}$
- IMPORTANT!!!! Here the number of items in $\mathcal{X}, m=|\mathcal{X}|$, is not fixed and may therefore vary.


## A closer look

- The probability mass function is fairly simple:

$$
p(\mathcal{X}=X) \propto \operatorname{det} \mathbf{L}_{X}
$$

- $\operatorname{det} \mathbf{L}_{\mathcal{X}} \geqslant 0$, by positive-definiteness of $\mathbf{L}$
- In addition: $\sum_{\mathcal{X}} \operatorname{det} \mathbf{L}_{\mathcal{X}}=\operatorname{det}(\mathbf{L}+\mathbf{I})$ is the normalisation constant (tractable!)
- So why does this induce repulsion?


## Determinants: geometric interpretation



Determinants measure the (signed) volume of the paralleliped spanned by the columns of a matrix. Illustration by Yigit Pilavci.

Why does the determinant induce repulsion?



Determinant: 0.51 .

Why does the determinant induce repulsion?


$\mathbf{L}_{\mathcal{X}}=$|  | 67 | 178 | 125 |
| ---: | ---: | ---: | ---: |
| 67 | 1.00 | 0.95 | 0.89 |
| 178 | 0.95 | 1.00 | 0.97 |
| 125 | 0.89 | 0.97 | 1.00 |

Determinant: 0.005 .

## Inclusion probabilities

- Are certain, or pairs of items are more likely to be sampled?
- Formally: let $\mathcal{S}$ denote a fixed (non-random) set. The "inclusion probabilities" are of the form:

$$
p(\mathcal{S} \subseteq \mathcal{X})
$$

- If $\mathcal{S}=\{i\}$, a singleton, equivalent to $p(i \in \mathcal{X})$, the probability that item $i$ is sampled
- If $\mathcal{S}=\{i, j\}$, a pair, equivalent to $p(i \in \mathcal{X}$ and $j \in \mathcal{X})$, the probability that both items are sampled


## Marginal kernels

- In DPPs the inclusion probabilities are quite remarkable
- For a DPP with L-ensemble $\mathbf{L}$ the inclusion probabilities are as follows

$$
p(\mathcal{S} \subseteq \mathcal{X})=\operatorname{det} \mathbf{K}_{\mathcal{S}}
$$

where:

$$
\mathbf{K}=\mathbf{L}(\mathbf{L}+\mathbf{I})^{-1}
$$

- $\mathbf{K}$ is called the marginal kernel of the DPP


## L-ensemble vs. marginal kernel

Example.

$$
\mathbf{L}=\left(\begin{array}{ccccc}
1 & 0.946 & 0.681 & 0.634 & 0.611 \\
0.946 & 1 & 0.864 & 0.825 & 0.805 \\
0.681 & 0.864 & 1 & 0.997 & 0.993 \\
0.634 & 0.825 & 0.997 & 1 & 0.999 \\
0.611 & 0.805 & 0.993 & 0.999 & 1
\end{array}\right)
$$

can be used to compute $p(\mathcal{X}=X)$.

$$
\mathbf{K}=\mathbf{L}(\mathbf{L}+\mathbf{I})^{-1}=\left(\begin{array}{lllll}
0.328 & 0.246 & 0.075 & 0.053 & 0.042 \\
0.246 & 0.234 & 0.135 & 0.117 & 0.108 \\
0.075 & 0.135 & 0.206 & 0.210 & 0.212 \\
0.053 & 0.117 & 0.210 & 0.219 & 0.223 \\
0.042 & 0.108 & 0.212 & 0.223 & 0.227
\end{array}\right)
$$

can be used to compute $p(\mathcal{S} \in \mathcal{X})$

## First-order inclusion probabilities

- First-order inclusion probabilities are just:

$$
p(i \in \mathcal{X})=K_{i i}
$$

- Exercise: work out $E(|\mathcal{X}|)$
- Hint: $|\mathcal{X}|=\sum_{j \in \Omega} \mathbb{I}(j \in \mathcal{X})$

First-order inclusion probabilities are (generally) not uniform!


Radius prop. to $p(i \in \mathcal{X})=K_{i i}$

## Second-order inclusion probabilities

- Note $\pi_{i}=p(i \in \mathcal{X})$
- Poisson sampling : go through all $n$ items and include item $i$ with probability $\pi_{i}$ independently
- Exercise: let $\mathcal{Y}$ be a Poisson sample with the same first-order inclusion probabilities as $\mathcal{X}$. Compute $p(i, j \subseteq \mathcal{Y})$. Compare to $p(i, j \subseteq \mathcal{X})$ : how does repulsion manifest itself?


## Fixed-size DPPs

- Often it's preferable to set the size of $\mathcal{X}$ to a fixed value.
- A fixed-size DPP is a DPP, conditioned on $|\mathcal{X}|=m$. They were introduced by Kulesza \& Taskar as "k-DPPs". Here we call them "m-DPPs" for consistency.
- Def. $\mathcal{X}$ is a m-DPP with L-ensemble $\mathbf{L}$ if

$$
p(\mathcal{X})= \begin{cases}\frac{\operatorname{det} \mathbf{L}_{\mathcal{X}}}{e_{m}(\mathbf{L})} & \text { if }|\mathcal{X}|=m \\ 0 & \text { otherwise }\end{cases}
$$

- $e_{m}(\mathrm{~L})$ is the normalisation constant, and is easy to compute from the spectrum of L.
- Otherwise an m-DPP is very similar to a DPP: we're simply forbidding sets of a size smaller or greater than $m$


## Inclusion probabilities in m-DPPs

- The bad news: m-DPPs do not, in general, have a marginal kernel, i.e. there may not be a matrix K such that

$$
p(\mathcal{S} \subseteq \mathcal{X})=\operatorname{det} \mathbf{K}_{\mathcal{S}}
$$

when $\mathcal{S}$ is a m-DPP.

- Exact inclusion probabilities are tricky to compute, especially for $|\mathcal{S}|>1$


## Inclusion probabilities in m-DPPs

- The good news: we showed in Barthelmé, Tremblay, Amblard (2019) that there is an approximate marginal kernel, i.e. for large $n$ and small $|\mathcal{S}|$, there's a matrix $\tilde{\mathbf{K}}$ such that

$$
p(\mathcal{S} \subseteq \mathcal{X}) \approx \operatorname{det} \tilde{\mathbf{K}}_{\mathcal{S}}
$$

- $\tilde{\mathbf{K}}$ is easy to compute:

$$
\tilde{\mathbf{K}}=\alpha \mathbf{L}(\alpha \mathbf{L}+\mathbf{I})^{-1}
$$

where $\alpha$ is such that $\operatorname{Tr} \tilde{\mathbf{K}}=m$

## Projection DPPs

- m-DPPs do not have exact marginal kernels, with one very important exception
- If $m=r=\operatorname{rank} \mathbf{L}$, then there is an exact marginal kernel, with a very specific form
- Let $\mathbf{L}=\mathbf{U D U}{ }^{t}$, the eigendecomposition of $\mathbf{L}$, and $\mathbf{D}$ the $r \times r$ matrix of eigenvalues.
- The marginal kernel is simply $\mathbf{K}=\mathbf{U} \mathbf{U}^{t}$, a projection matrix $\left(\mathbf{K}^{2}=\mathbf{K}\right)$
- Accordingly these DPPs are called projection DPPs.
- In a sense they are both DPPs and m-DPPs
- They are central to the overall theory


## An example of a projection DPP

- Here's an example of how to build a projection DPP. Assume the items are just points along a line: $x_{1}, \ldots, x_{n}$.
- We build a matrix of polynomial features:

$$
\mathbf{M}=\left(\begin{array}{ccccc}
1 & x_{1} & x_{1}^{2} & \ldots & x_{1}^{r-1} \\
1 & x_{2} & x_{2}^{2} & \ldots & x_{2}^{r-1} \\
\vdots & \vdots & \vdots & & \vdots \\
1 & x_{n} & x_{n}^{2} & \ldots & x_{n}^{r-1}
\end{array}\right)
$$

- We build an L-ensemble based on those features:

$$
\mathbf{L}=\mathbf{M} \mathbf{M}^{t}
$$

- L has rank $r$ and dimension $n \times n$
- If we set $m=r$, ie. we sample as many points as we have polynomial features, than what we have is a projection DPPs.


## Summary so far

- DPPs have tractable inclusion probabilities, but the number of items sampled is random (in general)
- m-DPPs have fixed sample size, but the inclusion probabilities are less tractable
- One exception: projection DPPs have fixed sample size, and the inclusion probabilities are tractable

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## Some computational issues

- There's a few computational issues, but we'll look at the two main ones:

1. How to sample from a DPP efficently
2. How to create an L-ensemble efficiently

- We can't cover the theory in detail so focus is on practical aspects
- See our package DPP.jl for efficient Julia implementation; DPPy by Guillaume Gautier for Python tools


## Samplers for DPPs

- For DPPs there are both exact and inexact samplers
- The inexact samplers (eg. Gibbs sampler) use an MCMC chain to generate approximate samples cheaply.
- However getting an exact sample is often not much more expensive: we will describe a method based on Hough et al. (2006)

A Metropolis-Hastings sampler

## Initial configuration

A Metropolis-Hastings sampler

## Propose swap

## A Metropolis-Hastings sampler

## Compute acceptance ratio



A Metropolis-Hastings sampler

## Accept or reject new configuration

## A Metropolis-Hastings sampler

Initialisation: set $\mathcal{X}$ to some random subset of size $m$. For $t=1$ to $T$, do:

- Propose swap: construct $\mathcal{X}^{\prime}$ by removing random item from $\mathcal{X}$, adding a random item from $\Omega-\mathcal{X}$
- Evaluate acceptance ratio $r=\frac{\operatorname{det} \mathbf{L}_{\mathcal{X}^{\prime}}}{\operatorname{det} \mathbf{L}_{\mathcal{X}}}$
- Set $\mathcal{X} \leftarrow \mathcal{X}^{\prime}$ with probability $r$.

If $T$ is large enough, the final configuration is an almost-exact sample from an m-DPP with ensemble $\mathbf{L}$

## A Metropolis-Hastings sampler

- The sampler we've just described is really easy to implement!
- Feel free to try it for yourself after the tutorial, should just take a few minutes
- Bonus points if you can adapt it to DPPs and not just m-DPPs
- For most practical purposes we recommend the exact sampler we describe next


## The direct sampler

- It turns out that sampling from a projection DPP is easy
- The algorithm just picks points sequentially from the appropriate probability distribution
- For generic DPPs, we'll see that it's possible to reduce the problem to the sampling of a projection DPP


## Sampling sequentially

- Take a set of $m$ items $\mathcal{X}=\left\{x_{1}, \ldots, x_{m}\right\}$ and order it (arbitrarily) into a sequence $x_{1}, \ldots, x_{m}$
- Our goal is to sample $\mathcal{X} \sim \operatorname{proj}-\operatorname{DPP}(K)$ by sampling first $x_{1}$, then $x_{2}$, then $x_{3}$ etc. up to a $x_{m}$
- Formally:

$$
\begin{array}{r}
x_{1} \sim p\left(x_{1}\right) \\
x_{2} \sim p\left(x_{2} \mid x_{1}\right) \\
x_{3} \sim p\left(x_{3} \mid x_{1}, x_{2}\right)
\end{array}
$$

## What are these conditional distributions?

- $p\left(x_{1}\right)$ is the distribution of an arbitrary item taken from a projection DPP - that's just the inclusion probability
- $p\left(x_{2} \mid x_{1}\right)$ is the distribution of an arbitrary item taken from a projection DPP, given that item $x_{1}$ is in the set. That's a conditional inclusion probability.
- etc.
- As it turns out, these distributions are tractable in proj-DPPs, and this leads to an algorithm that is both easy to implement and fast ${ }^{2}$
- Nice bit of theory: conditional distribution of $x_{t}$ equals the conditional variance of a Gaussian process with the same kernel sampled at $x_{1} \ldots x_{t-1}$ !

[^11]The direct sampling algorithm in action


The direct sampling algorithm in action


The direct sampling algorithm in action


The direct sampling algorithm in action


The direct sampling algorithm in action


## Sampling generic DPPs

- At this point we know how to sample from projection DPPs
- Now we need to sample from generic (m)-DPPs
- Luckily, we can show that generic (m)-DPPs are just mixtures of projection DPPs
- Recall: to sample from a mixture of Gaussians, pick randomly one of the Gaussians then sample from that Gaussian.
- Same here: we'll have to form a random projection DPP, then sample from that projection DPP


## Cauchy-Binet lemma

- We'll sketch the proof that all DPPs are mixtures of projection DPPs.
- Central ingredient is the Cauchy-Binet lemma.
- Let $\mathbf{A}_{m \times n}$ and $\mathbf{B}_{n \times m}$, with $n \geqslant m$. We seek to compute $\operatorname{det} \mathbf{A B}$.
- If $n=m \mathbf{A}$ and $\mathbf{B}$ are square, and so $\operatorname{det} \mathbf{A B}=\operatorname{det} \mathbf{A} \operatorname{det} \mathbf{B}$. Cauchy-Binet generalises this formula to $n>m$.

$$
\operatorname{det} \mathbf{A B}=\sum_{|\mathcal{Y}|=m} \operatorname{det} \mathbf{A}_{:, \mathcal{Y}} \operatorname{det} \mathbf{B}_{\mathcal{Y},:}
$$

- Here $\mathcal{Y}$ is a subset of $1,2, \ldots, n$ of size $m$ and the sum runs over all such subsets.


## Proof sketch that DPPs are mixtures of projection DPPs

Consider the eigendecomposition of $\mathbf{L}, \mathbf{L}=\mathbf{U D} \mathbf{U}^{t}$, and the probability of set $\mathcal{X}$.

$$
p(\mathcal{X}) \propto \operatorname{det} \mathbf{L}_{\mathcal{X}}=\operatorname{det} \mathbf{U}_{\mathcal{X},:} \mathbf{D}^{\frac{1}{2}} \mathbf{D}^{\frac{1}{2}} \mathbf{U}_{:, \mathcal{X}}^{t}=(\operatorname{det} \mathbf{A B})
$$

Apply Cauchy-Binet:

$$
p(\mathcal{X}) \propto \sum_{|\mathcal{Y}|=|\mathcal{X}|} \operatorname{det} \mathbf{U}_{\mathcal{X}, \mathcal{Y}} \mathbf{D}_{\mathcal{Y}, \mathcal{Y}}^{\frac{1}{2}} \operatorname{det} \mathbf{D}_{\mathcal{Y}, \mathcal{Y}}^{\frac{1}{2}} \mathbf{U}_{\mathcal{Y}, \mathcal{X}}^{t}
$$

Now we have square matrices inside the sum, so this is just:

$$
p(\mathcal{X}) \propto \sum_{|\mathcal{Y}|=|\mathcal{X}|} \operatorname{det} \mathbf{U}_{\mathcal{X}, \mathcal{Y}} \mathbf{U}_{\mathcal{Y}, \mathcal{X}}^{t} \operatorname{det} \mathbf{D}_{\mathcal{Y}, \mathcal{Y}}
$$

## Proof sketch that DPPs are mixtures of projection DPPs

Looking at:

$$
p(\mathcal{X}) \propto \sum_{|\mathcal{Y}|=|\mathcal{X}|} \operatorname{det} \mathbf{U}_{\mathcal{X}, \mathcal{Y}} \mathbf{U}_{\mathcal{Y}, \mathcal{X}}^{t} \operatorname{det} \mathbf{D}_{\mathcal{Y}, \mathcal{Y}}
$$

with $\mathcal{X}$ as a variable, we see the following structure appearing:

$$
p(\mathcal{X}) \propto \sum_{|\mathcal{Y}|=|\mathcal{X}|} f(\mathcal{X} \mid \mathcal{Y}) g(\mathcal{Y})
$$

which expresses $p(\mathcal{X})$ as a marginal! Here $f(\mathcal{X} \mid \mathcal{Y})=\operatorname{det} \mathbf{U}_{\mathcal{X}, \mathcal{Y}} \mathbf{U}_{\mathcal{Y}, \mathcal{X}}^{t}$, and that's a projection DPP where we select the eigenvectors given by $\mathcal{Y}$ to form the L-ensemble. $g(\mathcal{Y})=\operatorname{det}\left(\mathbf{D}_{\mathcal{Y}}\right)$ is also a DPP, this time with a diagonal L-ensemble!

## Computational considerations

- A tally of computational costs:

1. We need to generate the $\mathbf{L}$ matrix at cost $\mathcal{O}\left(n^{2}\right)$
2. We need to compute the eigendecomposition of $\mathbf{L}$ at $\operatorname{cost} \mathcal{O}\left(n^{3}\right)$
3. We need to sample at cost $\mathcal{O}\left(n k^{2}\right)$

## Computational considerations

- Overall the dominating cost is the eigendecomposition at cost $\mathcal{O}\left(n^{3}\right)$
- Fortunately that cost can be brought down to $\mathcal{O}\left(n k^{2}\right)$ if you design $\mathbf{L}$ to have rank $\mathcal{O}(k)$
- For example: use $k$ polynomial features, or use Random Fourier Features
- See Tremblay et al. (2018) on coresets ${ }^{3}$ for a list of tricks
- In a nutshell: DPPs have very good (linear) scaling in $n$, meaning the original set can be in the millions
- However, poor scaling in $k$, so that the subset you sample will be in the hundreds

[^12]Introduction
DPPs to produce diverse samples
DPPs as a tool in SP/ML
DPPs to characterize

Definition, basic properties
Repulsive point processes are hard DPPs, the nitty-gritty

Computation
Sampling from a DPP
DPPs as mixtures

Applications
Examples of applications
Zoom on an application: Coresets

Conclusion

## Example of application: generate extractive summaries ${ }^{1}$

```
NASA and the Russian Space Agency have agreed to set
    mside s last-minute Russian request to lsunch an
    internationa< space station into an orbit closer to Mir,
    ofticials announced Friday.
```

    A last-minute alarm forced NASA to halt Thursday's
    launching of the space shuttie Endeavour, on a mission to
    start assembling the intemational space station. This wat
    the first time in three years.
    The plaret s most diring construction job begar Friday as
    the shuttie Endeavour carried into orbie sue astronauts and
    the first US.-buit part of an internaticenal space station
    that is espected ta cost more than \(\$ 100\) bilfion.
    Following a series of intricate maneuvers and the skiliful
    use of the space shuttle Endeavour'\& rabot arm.
    astronauts on Sunday joined the first two of mary
    segments that will form the space station.
    ***
    On Fridyy the shuttie Indeavor carried six astronauts into orbet to start
        buiding an international space station The launch occurred after fussa
        and U.5. officials agreed not to delpy the flight in order to orbit cioser to
    NR, and after a last-minute aiserm forced a pontponement. On Suandey
    astronauts foining the Russ an-made Zarya control modute cylnder with
    the American-made module to form a 70.000 pounds mass 77 feet.
    tong.
    human summary
    - NASA and the Russian Space Agency have agreed to set aside
- A last-minute alarm forced NASA to halt Thursdays launching.
- This was the first time in three yeark and 19 fights.
+After a last-minute alarm, the launch wert off liawlensly Fridsy
+ Following a series of intricate maneuvers and the skllith
- It looked to be a perfect and, hopefully, long-lasting fit.
extractive summary
document cluster
- The trick is to find a good feature space to embed sentences and a proper DPP kernel
- Both can be parametrized (tf, idf, etc.) and then learned ${ }^{2}$

[^13]DPP Tutorial

## Example of application: search algorithms ${ }^{1}$

"porsche"
$k=2$

"philadelphia"


[^14]
## "DDPs as a tool" applications

- Monte-Carlo integration ${ }^{12}$ :

$$
\int f(x) \mu(d x) \simeq \sum_{n=1}^{N} \omega_{n} f\left(x_{n}\right)
$$

where the $x_{i}$ 's are the so-called quadratic nodes.

- Mini-batch sampling for stochastic gradient descent ${ }^{3}$ :

$$
\begin{aligned}
& L(\theta)=\sum_{i} L_{i}(\theta) \\
& \mathrm{GD}: \theta \leftarrow \theta-\eta \nabla L(\theta)=\theta-\eta \sum_{i} \nabla L_{i}(\theta) \\
& \text { mini-batch GD }: \theta \leftarrow \theta-\eta \sum_{i \in \mathcal{X}} \nabla L_{i}(\theta)
\end{aligned}
$$

- Column subset selection problem for best rank- $k$ approximation ${ }^{4}$

[^15]
## Zoom on one application:

- Coresets ${ }^{1}$
${ }^{1}$ Tremblay et al., DPPs for Coresets, Arxiv, 2018.


## Coresets

- Consider a dataset $\mathcal{X}=\left(x_{1}, \ldots, x_{n}\right)$, say: $n$ points in dimension $d$.
- Let $\Theta$ be a parameter space and consider cost functions of the form:

$$
L(\mathcal{X}, \theta)=\sum_{i=1}^{n} f\left(\boldsymbol{x}_{i}, \theta\right)
$$

where $f: \mathcal{X} \rightarrow \mathbb{R}^{+}$, and $\theta \in \Theta$.

- A classical ML objective: find

$$
\theta^{*}=\underset{\theta \in \Theta}{\operatorname{argmin}} L(\mathcal{X}, \theta) .
$$

- $k$-means, $k$-medians, linear/logistic regressions fall in this class of problems


## Coresets

- Consider a subset $\mathcal{S} \subset \mathcal{X}$ (possibly with repetitions)
- Associate a weight $\omega_{s}>0$ to each element $\boldsymbol{s} \in \mathcal{S}$
- Define

$$
\hat{L}(\mathcal{S}, \theta)=\sum_{s \in \mathcal{S}} \omega_{s} f(\boldsymbol{s}, \theta)
$$

## Coresets

- $\mathcal{S}$ is an $\epsilon$-coreset of $\mathcal{X}$ wrt $L$ if:

$$
\forall \theta \in \Theta \quad(1-\epsilon) L(\mathcal{X}, \theta) \leqslant \hat{L}(\mathcal{S}, \theta) \leqslant(1+\epsilon) L(\mathcal{X}, \theta)
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- Multiplicative approximation: gold standard of approximation methods


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\hat{L}\left(\mathcal{S}, \hat{\theta}^{*}\right)
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## Coresets: illustration on the 1-means problem

- Data $\mathcal{X}$



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L(\mathcal{X}, \theta)=\sum_{i=1}^{n}\left\|\boldsymbol{x}_{i}-\theta\right\|^{2}
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$$
\forall \theta \quad\left|\frac{\hat{L}}{L}-1\right| \leqslant \epsilon
$$

- Estimated optimal $\theta$ :

$$
\hat{\theta}^{*}=\underset{\theta \in \Theta}{\operatorname{argmin}} \hat{L}(\mathcal{S}, \theta)
$$




## Random coresets

- Random context: suppose $\mathcal{S}$ is a random subset $\mathcal{S} \subset \mathcal{X}$ (possibly with repetitions)
- Importance sampling notations:
- Define $\epsilon_{i}$ the random variable counting the number of times $\boldsymbol{x}_{i}$ is in $\mathcal{S}$
- To each element $x_{i}$ associate a weight $\omega_{i}=\frac{1}{\mathbb{E}\left(\epsilon_{i}\right)}$
- One has:

$$
\hat{L}(\mathcal{S}, \theta)=\sum_{i=1}^{n} f\left(\boldsymbol{x}_{i}, \theta\right) \frac{\epsilon_{i}}{\mathbb{E}\left(\epsilon_{i}\right)}
$$

and thus $\hat{L}$ is an unbiased estimator of $L$ :

$$
\mathbb{E}(\hat{L}(\mathcal{S}, \theta))=\sum_{i=1}^{n} f\left(\boldsymbol{x}_{i}, \theta\right)=L(\mathcal{X}, \theta) .
$$

## Sensitivity

- The sensitivity of a datapoint $x_{i} \in \mathcal{X}$ with respect to $f: \mathcal{X}, \Theta \rightarrow \mathbb{R}^{+}$is:

$$
\sigma_{i}=\max _{\theta \in \Theta} \frac{f\left(\boldsymbol{x}_{i}, \theta\right)}{L(\mathcal{X}, \theta)} \quad \in[0,1] .
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- 1-means is an exception. In this case, supposing wlog that the data is centered (i.e.: $\sum_{j} x_{j}=0$ ), one shows:

$$
\sigma_{i}=\frac{1}{n}\left(1+\frac{\left\|x_{i}\right\|^{2}}{v}\right)
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where $v=\frac{1}{n} \sum_{x \in \mathcal{X}}\|x\|^{2}$. Thus, $\mathfrak{S}=2$.

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## A classical iid coreset theorem ${ }^{1}$

- Let $\boldsymbol{p} \in[0,1]^{n}$ be a probability distribution over all datapoints $\mathcal{X}$ with $p_{i}$ the probability of sampling $x_{i}$ and $\sum_{i} p_{i}=1$.

[^16]
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- Draw $\mathcal{S}$ : $m$ iid samples with replacement according to $\boldsymbol{p}$.


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- Associate importance sampling weights to each sample of $\mathcal{S}$.

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- Draw $\mathcal{S}: m$ iid samples with replacement according to $\boldsymbol{p}$.
- Associate importance sampling weights to each sample of $\mathcal{S}$.
- Theorem. The weighted subset $\mathcal{S}$ is a $\epsilon$-coreset with high probability if:

$$
m \geqslant \mathcal{O}\left(\frac{d^{\prime}}{\epsilon^{2}}\left(\max _{i} \frac{\sigma_{i}}{p_{i}}\right)^{2}\right)
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where $d^{\prime}$ is the pseudo-dimension of $\Theta$ (a generalization of the Vapnik-Chervonenkis dimension).

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- The optimal probability distribution minimizing the rhs is $p_{i}=\sigma_{i} / \mathfrak{G}$.
- In this case, $\mathcal{S}$ is a $\epsilon$-coreset with high probability if:

$$
m \geqslant \mathcal{O}\left(\frac{d^{\prime} \mathfrak{S}^{2}}{\epsilon^{2}}\right)
$$

[^20]
## DPPs for Coresets: a result ${ }^{1}$

- Consider any iid sampling scheme, defined by:
- $m$ the number of samples to draw
- $\forall i, 0 \leqslant p_{i} \leqslant 1 / m$ and $\sum_{i} p_{i}=1$


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- Consider a marginal kernel K verifying:
- K is projective of rank $m: \mathrm{K}=U U^{t}$ with $U \in \mathbb{R}^{n \times m}$ and $U^{t} U=I_{m}$.
- $\forall i, \mathrm{~K}_{i i}=m p_{i}$.


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- Sample $\mathcal{S}_{\text {iid }}$ by drawing $m$ samples iid from $\boldsymbol{p}$
- Sample $\mathcal{S}_{d p p}$ from the DPP of kernel K.
- Recall that $\mathcal{S}_{d p p}$ is necessarily of size $m$.

[^22]
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- Sample $\mathcal{S}_{\text {iid }}$ by drawing $m$ samples iid from $\boldsymbol{p}$
- Sample $\mathcal{S}_{d p p}$ from the DPP of kernel K.
- Recall that $\mathcal{S}_{d p p}$ is necessarily of size $m$.
- Coreset variance reduction theorem. One has:

$$
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$\rightarrow$ We propose a computationally efficient heuristic based on the Gaussian kernel:
- Compute $r$ Random Fourier Features ${ }^{2}(r=\mathcal{O}(m))$ and obtain $\psi \in \mathbb{R}^{n \times r}$ s.t. $\Psi \Psi^{t} \in \mathbb{R}^{n \times n}$ approximates the Gaussian kernel
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$\rightarrow$ This runs in $\mathcal{O}\left(n m^{2}+n m d\right)$

[^30]In practice: the 1-means controlled example ${ }^{1}$

- Data $\mathcal{X}$, parameter $\theta$
- Cost func.

$$
L(\mathcal{X}, \theta)=\sum_{i=1}^{n}\left\|\boldsymbol{x}_{i}-\theta\right\|^{2}
$$

## Compare:

- uniform iid sampling
- sensitivity iid: ideal iid sampling based on exact sensitivities
- m-DPP (heuristic) based on RFFs of the Gaussian L-ensemble

$$
\mathrm{L}_{i j}=\exp ^{-\left\|x_{i}-x_{j}\right\|^{2} / s^{2}}
$$



[^31]
## Conclusion: take home messages

- DPPs create random, diverse samples.
- They are tractable (inclusion probabilities at all orders are known), and good approximations are known for m-DPPs.
- This does not mean they are the best choice for all applications! They are many other (less tractable) repulsive processes out there.
- They are used in practice
- They are not expensive to sample in many applications (where low-rank approximations of the kernel can be computed efficiently): $\mathcal{O}\left(n m^{2}\right)$
- Toolboxes exist: DPPy ${ }^{1}$ in Python, DPP.j1 ${ }^{2}$ in Julia

[^32]
## Conclusion

DPPs have links with many other theories:

- graph theory
- Gaussian processes
- multivariate polynomials
- random matrices
- etc.


## Conclusion: what next?

- accelerate sampling for large $m$
- DPPs for large dimensional data?
- parallel implementations


[^0]:    ${ }^{1}$ left figure: from Kulesza and Taskar, DPPs for machine learning, Found. and Trends in ML, 2013
    ${ }^{2}$ right figure: from G. Gautier's slides guilgautier.github.io/pdfs/GaBaVa17_slides.pdf
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[^2]:    ${ }^{1}$ see, e.g., Johansson, Random matrices and DPPs, Arxiv (lecture notes), 2005

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[^6]:    ${ }^{1}$ see, e.g., Johansson, Random matrices and DPPs, Arxiv (lecture notes), 2005

[^7]:    ${ }^{1}$ Macchi, The coincidence approach to stochastic point processes. Adv. Appl. Probab., 1975

[^8]:    ${ }^{1}$ Macchi, The coincidence approach to stochastic point processes. Adv. Appl. Probab., 1975

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[^10]:    ${ }^{1}$ Löwen, H. (2000). Fun with hard spheres. In Statistical physics and spatial statistics (pp. 295-331). Springer, Berlin, Heidelberg.

[^11]:    ${ }^{2}$ Alg. due to Hough et al. (2006), Gillenwater (2014) for a faster version. See DDPy documentation by G. Gautier for more

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[^13]:    ${ }^{2}$ e.g., Kulesza et al, Near-optimal map inference for DPPs, NIPS 2012.
    ${ }^{1}$ Kulesza and Taskar, DPPs for machine learning, Found. and Trends in ML, 2013

[^14]:    ${ }^{1}$ Kulesza and Taskar, DPPs for machine learning, Found. and Trends in ML, 2013

[^15]:    ${ }^{1}$ Gautier et al., On two ways to use DPPs for Monte Carlo integration, ICML, 2019.
    ${ }^{2}$ Bardenet et al., Monte Carlo with DPPs, Annals of Applied Probability, In Press.
    ${ }^{3}$ Zhang et al., DPPs for Mini-Batch Diversification, UAI, 2017.
    ${ }^{4}$ Belhadji et al, A DPP for column subset selection, Arxiv, 2018.

[^16]:    ${ }^{1}$ Langberg and Schulman, Universal $\epsilon$-approximators for integrals, SIAM, 2010

[^17]:    ${ }^{1}$ Langberg and Schulman, Universal $\epsilon$-approximators for integrals, SIAM, 2010

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[^21]:    ${ }^{1}$ Tremblay et al., DPPs for Coresets, Arxiv, 2018.

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[^24]:    ${ }^{2}$ Rahimi et al., Random features for large-scale kernel machines, NIPS, 2008

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[^31]:    ${ }^{1}$ Tremblay et al., DPPs for Coresets, Arxiv, 2018.

[^32]:    ${ }^{1}$ github.com/guilgautier/DPPy
    ${ }^{2}$ gricad-gitlab.univ-grenoble-alpes.fr/barthesi/dpp.jl

